

# on the duality between disks and trees

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## background

André Joyal defined the category *Disk* and implicitly its opposite category  $\Theta$  in order to define weak higher categories.

An explicit definition of  $\Theta$  in terms of trees and (strict)  $\omega$ -categories was provided by Michael Batanin and Ross Street.

Michael Makkai and Marek Zawadowski and then Clemens Berger provided demonstrations of the duality between *Disk* and  $\Theta$ .

## overview

Here we demonstrate the duality between *Disk* and  $\Theta$  by lifting the well known duality between ordinals and intervals.

We do so by extending the category of ordinals and the category of intervals using induction and also using so-called labeled trees.

The extended categories of intervals are equivalent to *Disk*.

The extended categories of ordinals are equivalent to  $\Theta$ .

# ordinals and intervals

$\Delta_+$  — the category of ordinals

- ▶ objects: finite ordinals  $[n] = \{0, \dots, n\}$  for  $n = -1, 0, \dots$
- ▶ morphisms: order preserving functions

$\mathcal{I}_+$  — the category of intervals

- ▶ objects: non-empty ordinals
- ▶ morphisms: greatest, least and order preserving functions

There is a well known equivalence  $\mathcal{I}_+^{op} \cong \Delta_+$  by the functors

$$\begin{aligned} (-)^\wedge : \Delta_+^{op} &\rightarrow \mathcal{I}_+ & (-)^\vee : \mathcal{I}_+^{op} &\rightarrow \Delta_+ \\ &: [n] \mapsto [n+1] & &: [n] \mapsto [n-1]. \end{aligned}$$

Next we introduce categories named  $i\mathcal{I}_+$  and  $i\Delta_+$  which are inductively defined extensions of  $\mathcal{I}_+$  and  $\Delta_+$ .

## the category $i\mathcal{I}_+$

$i\mathcal{I}_+$  — an extension of the intervals defined inductively

The objects of  $i\mathcal{I}_+$  are graded by height.

The interval  $[0]$  is the trivial object of  $i\mathcal{I}_+$  and has height 0 (zero).

$H$ , an object of height  $n$ , consists of

- ▶ an interval  $\text{Ob } H$ , called the root object of  $H$
- ▶ an object  $H(i)$  of  $i\mathcal{I}_+$  of height  $< n$  for each  $i \in \text{Ob } H$  which is trivial if and only if  $i$  is an endpoint

A morphism  $g: H \rightarrow H'$  consists of

- ▶ an interval morphism  $g: \text{Ob } H \rightarrow \text{Ob } H'$
- ▶ a morphism  $g(i): H(i) \rightarrow H'(gi)$  for each  $i \in \text{Ob } H$

## the category $i\Delta_+$

$i\Delta_+$  — an extension of the ordinals defined inductively

The objects of  $i\Delta_+$  are graded by height.

The ordinal  $[-1]$  is the trivial object of  $i\Delta_+$  and has height 0 (zero).

$K$ , an object of height  $n$ , consists of

- ▶ an ordinal  $\text{Ob } K$ , called the root object of  $K$
- ▶ an object  $K(j)$  of  $i\Delta_+$  of height  $< n$  for each  $j \in (\text{Ob } K)^\wedge$  which is trivial if and only if  $j$  is an endpoint

A morphism  $g: K' \rightarrow K$  consists of

- ▶ an ordinal morphism  $g: \text{Ob } K' \rightarrow \text{Ob } K$
- ▶ a morphism  $g(j): K'(g^\wedge j) \rightarrow K(j)$  for each  $j \in (\text{Ob } K)^\wedge$

## equivalence between $i\mathcal{I}_+^{op}$ and $i\Delta_+$

The isomorphism

$$\mathcal{I}_+^{op} \cong \Delta_+$$

lifts to an isomorphism

$$i\mathcal{I}_+^{op} \cong i\Delta_+$$

using induction on the height of objects.

Next we

- ▶ define the category of trees
- ▶ recall the definition of Joyal's category *Disk*
- ▶ sketch the constructions used to demonstrate  $Disk \simeq i\mathcal{I}$

## the category of trees

A tree  $(A, p)$  is a functor  $A: \omega^{op} \rightarrow \text{Set}$

$$\dots \xrightarrow{p_2} A_2 \xrightarrow{p_1} A_1 \xrightarrow{p_0} A_0$$

where  $A_0$  is a singleton set containing the root of  $A$ .

A tree  $(A, p)$  has degree  $n$  when  $A_m \cong A_n$  for all  $m \geq n$ .

The tree over  $i$  for  $i \in A_n$  is the largest subtree with root  $i$ .

A morphism  $g: A \rightarrow B$  of trees is a natural transformation

$$\begin{array}{ccccccc} \dots & \xrightarrow{p_2} & \mathbf{A}_2 & \xrightarrow{p_1} & \mathbf{A}_1 & \xrightarrow{p_0} & \mathbf{A}_0 \\ & & \downarrow g_2 & & \downarrow g_1 & & \downarrow g_0 \\ \dots & & & & & & \\ \dots & \xrightarrow{q_2} & \mathbf{B}_2 & \xrightarrow{q_1} & \mathbf{B}_1 & \xrightarrow{q_0} & \mathbf{B}_0. \end{array}$$



## the categories $Disk$ and $Disk_+$

A disk as defined by Joyal is a tree  $(A, p)$  of finite degree

1. such that the fibers of  $p_n: A_{n+1} \rightarrow A_n$  have interval structure
2. with sections  $d_0, d_1: A_n \rightarrow A_{n+1}$  of  $p_n$   
where  $p_n^*(x) = [d_0(x), \dots, d_1(x)]$
3. such that the equalizer of  $d_0, d_1: A_n \rightarrow A_{n+1}$   
is  $d_0(A_{n-1}) \cup d_1(A_{n-1})$

A disk morphism is a tree morphism with components that are interval morphisms.

$Disk_+$  includes, in addition, the trees of degree 0 (zero).  
Such a tree is terminal in  $Disk_+$ .

## equivalence between $Disk_+$ and $i\mathcal{I}_+$

Given a disk  $(A, p)$  we construct inductively an object  $H$  of  $i\mathcal{I}_+$

- ▶  $\text{Ob } H$  is the interval fiber  $p_0^*(x)$  over the root of  $A$
- ▶ for every  $i \in \text{Ob } H$  let  $H(i)$  be the object of  $i\mathcal{I}_+$  constructed from the disk over  $i$

Given a disk morphism  $f: (A, p) \rightarrow (B, q)$

we construct inductively a morphism  $g: H \rightarrow K$  of  $i\mathcal{I}_+$

- ▶  $g: \text{Ob } H \rightarrow \text{Ob } K$  is the interval morphism  $p_0^*(x) \rightarrow q_0^*(y)$  between the fibers over the roots of  $A$  and  $B$
- ▶ for every  $i \in \text{Ob } H$  then  $g(i)$  is the morphism of  $i\mathcal{I}_+$  constructed from the restriction of  $f$  to the disk over  $i$

We have equivalences  $Disk_+ \simeq i\mathcal{I}_+$  and  $Disk \simeq i\mathcal{I}$   
that are surjective on objects.

## objects of $\Theta$

Batanin and Street defined  $\Theta$  to be the full subcategory of  $\omega\text{Cat}$  consisting of the free  $\omega$ -categories on globular cardinals.

We will soon define so-called ordinal graphs, which are inductively defined counterparts to globular cardinals, and use a similar method to construct objects of  $\Theta$ .

To construct an object of  $\Theta$  from an object of  $i\Delta_+$  we define

- ▶ the category of ordinal graphs
- ▶ the free  $\omega$ -category of an ordinal graph
- ▶ a map from objects of  $i\Delta_+$  to ordinal graphs

but first we recall the definition of  $\omega$ -category.

## $\omega$ -categories

an enriched category  $\mathcal{C}$  over a monoidal category  $\mathcal{V}$  is

- ▶ a set of objects  $\text{Ob } \mathcal{C}$
- ▶ a hom-object  $\mathcal{C}(A, B)$  in  $\mathcal{V}$  for every pair  $A, B \in \text{Ob } \mathcal{C}$
- ▶ with additional morphisms satisfying certain requirements

strict  $n$ -categories

- ▶ a 1-category is an ordinary category and is enriched over  $\text{Set}$
- ▶ an  $n$ -category is a category enriched over  $(n - 1)$ -categories

$\omega\text{Cat}$ , the category of  $\omega$ -categories, is the colimit of the diagram

$$0\text{-Cat} \rightarrow 1\text{-Cat} \rightarrow 2\text{-Cat} \rightarrow \cdots \rightarrow n\text{-Cat} \rightarrow \cdots$$

of inclusions.

# the category of ordinal graphs

*OGraph* — the category of  $\omega$ -graphs defined inductively

The objects of *OGraph* are graded by dimension.

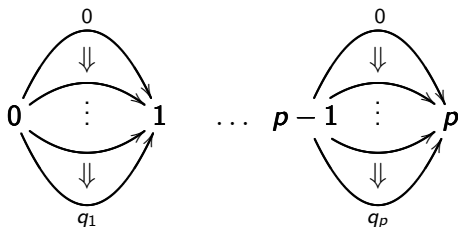
The ordinal graph of dimension -1 (minus one) is the empty set.

An ordinal graph  $\mathcal{G}$  of dimension  $n$  consists of

- ▶ a finite linearly ordered set  $\text{Ob } \mathcal{G} = \{x_0, \dots, x_p\}$  of vertices
- ▶ a non-empty ordinal graph  $\mathcal{G}(x_{i-1}, x_i)$  (an edge-object of  $\mathcal{G}$ ) of dimension  $< n$  for all  $i \in \{1, \dots, p\}$
- ▶  $\mathcal{G}(x_i, x_j)$  is empty when  $j \neq i + 1$

# ordinal graphs and globular cardinals

The ordinal graph  $[p]([q_1], \dots, [q_p])$  is drawn



This ordinal graph corresponds to a globular cardinal.  
In fact, we have an equivalence between

- ▶ the category of ordinal graphs
- ▶ the category of globular cardinals

## free $\omega$ -categories

The free  $\omega$ -category  $\mathfrak{F}\mathcal{G}$  on an ordinal graph  $\mathcal{G}$  has

- ▶ objects: those of  $\mathcal{G}$ , i.e.  $x_0, \dots, x_p$
- ▶ hom-objects: defined by induction

$$(\mathfrak{F}\mathcal{G})(x_i, x_j) = \mathfrak{F}(\mathcal{G}(x_{j-1}, x_j)) \times \dots \times \mathfrak{F}(\mathcal{G}(x_i, x_{i+1}))$$

The free  $\omega$ -category on an ordinal graph and the free  $\omega$ -category on the corresponding globular cardinal are isomorphic.

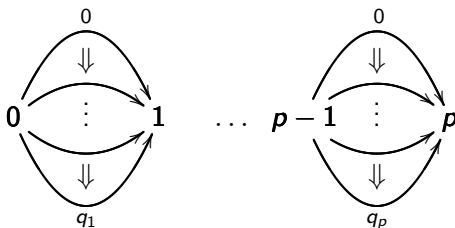
Objects of  $\Delta$  are free categories on ordinal graphs of dimension 1.

## from objects of $i\Delta_+$ to ordinal graphs

An object of  $i\Delta_+$  of height 2 (two) can be written

$$[p]([-1], [q_1], \dots, [q_p], [-1])$$

and corresponds to the ordinal graph



which is written

$$[p]([q_1], \dots, [q_p])$$

and we have a map  $\Upsilon: \text{Ob } i\Delta_+ \rightarrow \text{OGraph}$ .



from  $i\Delta_+$  to  $\omega\text{Cat}$  (on objects)

We have a map

$$\text{Ob } i\Delta_+ \xrightarrow{\Upsilon} \text{OGraph} \xrightarrow{\mathfrak{F}} \omega\text{Cat}$$

which sends  $H$ , an object of  $i\Delta_+$ , to the  $\omega$ -category  $\mathcal{C} = \mathcal{F}(H)$

- ▶ objects: elements of  $\text{Ob } H$  (an ordinal)
- ▶ hom-categories:  $\mathcal{C}(i, j)$  is the  $\omega$ -category

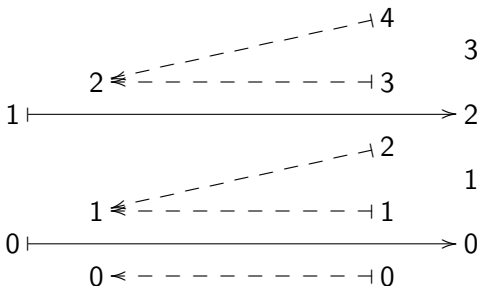
$$\prod_{\ell=i+1}^j \mathcal{F}\Upsilon H(\ell)$$

Next we construct an  $\omega$ -functor  $G: \mathcal{F}(K') \rightarrow \mathcal{F}(K)$   
given a morphism  $g: K' \rightarrow K$  of  $i\Delta_+$ .

# the ordinal/interval duality on morphisms

$g: K' \rightarrow K$  is a morphism of  $i\Delta_+$

- ▶ its object (ordinal) map  $g: [1] \rightarrow [3]$  is drawn with  $\rightarrow$
- ▶ the dual interval map  $g^\wedge: [4] \rightarrow [2]$  is drawn with  $\leftarrow$



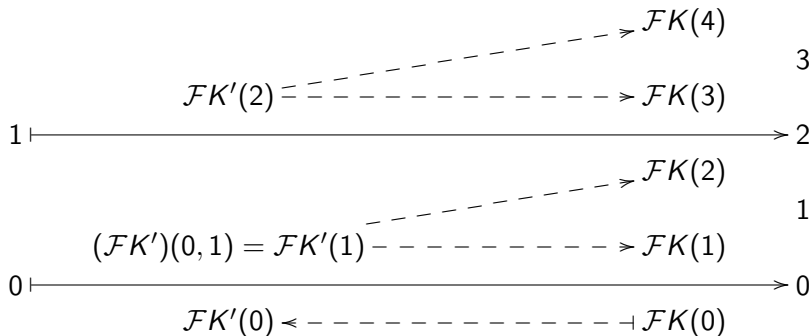
In general, we have  $g(j): K'(g^\wedge j) \rightarrow K(j)$  for each  $j \in (\text{Ob } K)^\wedge$ .  
In particular, we have the two morphisms of  $i\Delta_+$

$$g(1): K'(1) \rightarrow K(1) \quad \text{and} \quad g(2): K'(1) \rightarrow K(2).$$

## from $i\Delta_+$ to $\Theta$ (on morphisms)

We construct an  $\omega$ -functor  $G: \mathcal{F}(K') \rightarrow \mathcal{F}(K)$  from  $g: K' \rightarrow K$

- ▶ the object map of  $G$  is drawn below with  $\rightarrow$
- ▶ by induction we have  $\omega$ -functors drawn with  $--\rightarrow$
- ▶ the  $\omega$ -functors out of  $(\mathcal{F}K')(0, 1)$  induce an  $\omega$ -functor into the product  $\omega$ -category  $\mathcal{F}K(2) \times \mathcal{F}K(3)$  which is the hom-category  $(\mathcal{F}K)(G0, G1)$



## equivalence between $i\Delta_+$ and $\Theta$

These constructions (and some technical details) demonstrate an equivalence between  $i\Delta_+$  and  $\Theta_+$  where  $\Theta_+$  also contains the empty  $\omega$ -category in addition to the objects of  $\Theta$ .

We have the following chain of equivalences

$$Disk^{op} \xrightarrow{\simeq} i\mathcal{I}^{op} \xrightarrow{\simeq} i\Delta \xrightarrow{\simeq} \Theta.$$

## part II

The remainder of the talk describes two categories, named  $t\Delta_+$  and  $t\mathcal{I}_+$ , constructed from so-called labeled trees, and their reduced counterparts  $t\Delta$  and  $t\mathcal{I}$  which are also equivalent to the categories *Disk* and  $\Theta$  (respectively).

The equivalence  $t\mathcal{I}_+^{op} \simeq t\Delta_+$  is an easy consequence of the duality between ordinals and intervals.

Next we define the concept of a labeled tree.

## labeled trees

A labeled tree  $(A, F)$  is a tree  $A$  with functors  $F_n: A_n \rightarrow \mathcal{C}$

- ▶ and so is an object of  $Fam(\mathcal{C})$  for each  $n \in \mathbb{N}$

A morphism  $(g, \alpha)$  of labeled trees is a tree morphism  $g: A \rightarrow B$

- ▶ with natural transformations

$$\begin{array}{ccc} A_n & \xrightarrow{g_n} & B_n \\ & \searrow F_n & \swarrow G_n \\ & \mathcal{C} & \end{array} \quad \begin{array}{c} \alpha_n \\ \Rightarrow \end{array}$$

and so is a morphism of  $Fam_{\Sigma}(\mathcal{C})$  for each  $n \in \mathbb{N}$

An op-morphism has natural transformations

$$\begin{array}{ccc} A_n & \xrightarrow{g_n} & B_n \\ & \searrow F_n & \swarrow G_n \\ & \mathcal{C} & \end{array} \quad \begin{array}{c} \alpha_n \\ \Leftarrow \end{array}$$

- ▶ and so is a morphism of  $Fam_{\Pi}(\mathcal{C}^{op})$  for each  $n \in \mathbb{N}$

## constrained labeled trees

A constrained tree  $(A, F, \mathcal{U})$  is a labeled tree  $(A, F)$

- ▶ with functors  $F_n: A_n \rightarrow \mathcal{C}$  and  $\mathcal{U}: \mathcal{C} \rightarrow \text{Set}$
- ▶ such that  $A_{n+1} \cong \text{el}(\mathcal{U}F_n)$

A constrained morphism  $(f, \alpha): (A, F, \mathcal{U}) \rightarrow (B, G, \mathcal{U})$  is a morphism of labeled trees such that the following commutes

$$\begin{array}{ccc} A_{n+1} & \xrightarrow{\cong} & \text{el}(\mathcal{U}F_n) \\ f_{n+1} \downarrow & & \downarrow \text{el}(\mathcal{U}\alpha_n) \\ B_{m+1} & \xrightarrow{\cong} & \text{el}(\mathcal{U}G_n) \end{array}$$

$\text{Con}(\mathcal{C}, \mathcal{U})$  is the category of trees

- ▶ labeled by functors into  $\mathcal{C}$
- ▶ constrained by the functor  $\mathcal{U}: \mathcal{C} \rightarrow \text{Set}$

## the category $Cat/_{Set}$ and functor $Con$

The 2-category  $Cat/_{Set}$  has

▶ objects:  $(\mathcal{C}, \mathcal{U})$  with  $\mathcal{U}: \mathcal{C} \rightarrow Set$

▶ 1-cells: functors  $F: \mathcal{C} \rightarrow \mathcal{A}$  such that

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{F} & \mathcal{A} \\ & \searrow \mathcal{U}_{\mathcal{C}} & \swarrow \mathcal{U}_{\mathcal{A}} \\ & = & \\ & \text{Set} & \end{array}$$

▶ 2-cells: natural transformations  $\alpha: F \Rightarrow G$

The 2-functor  $Con: Cat/_{Set} \rightarrow Cat$  has

▶  $Con(\mathcal{C}, \mathcal{U})$  the category of trees labeled in  $\mathcal{C}$  constrained by  $\mathcal{U}$

▶  $Con(F)$  given by post-composition;  $Con(F)(A, H) = (A, FH)$

▶  $Con(\alpha)$  given by post-composition;  $Con(\gamma)(A, H) = (1_A, \gamma_H)$

$Con$  preserves equivalences and isomorphisms as it is a 2-functor.



## using the ordinal/interval duality

From the ordinal/interval duality we have this duality in  $Cat/Set$

$$(\mathcal{I}_+, U) \cong (\Delta_+, U(-)^\wedge)$$

and by the functor  $Con$  have the following duality in  $Cat$

$$Con(\mathcal{I}_+, U) \cong Con(\Delta_+, U(-)^\wedge)$$

where  $U$  is the ordinary ordinary underlying set functor.

Next we define full subcategories of the latter two categories which are equivalent to the categories  $Disk$  and  $\Theta$  (respectively).

## cropped trees

Let  $(A, F)$  be a tree labeled in  $\mathcal{I}_+$  constrained by  $U$

- ▶  $x \in A_{n+1}$  is an end element when it corresponds under the isomorphism  $A_{n+1} \cong \text{el}(UF_n)$  to an endpoint of  $F_n(y)$  for some  $y \in A_n$
- ▶  $(A, F)$  is cropped when  $F_n(x) = [0]$  for every end element  $x$

Let  $(A, F)$  be a tree labeled in  $\Delta_+$  constrained by  $U(\_)^{\wedge}$

- ▶  $x \in A_{n+1}$  is an end element when it corresponds under the isomorphism  $A_{n+1} \cong \text{el}U(F_n\_)^{\wedge}$  to an endpoint of  $(F_n y)^{\wedge}$  for some  $y \in A_n$
- ▶  $(A, F)$  is cropped when  $F_n(x) = [-1]$  for every end element  $x$

$t\mathcal{I}_+$  and  $t\Delta_+$

Define

- ▶  $t\mathcal{I}$  as the full subcategory of  $\text{Con}(\mathcal{I}_+, U)$   
whose objects are cropped trees of positive degree
- ▶  $t\Delta$  as the full subcategory of  $\text{Con}(\Delta_+, U(\_)^\wedge)$   
whose objects are cropped trees of positive degree

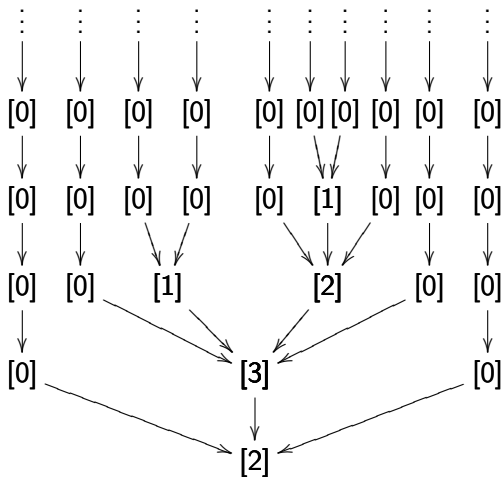
We have an equivalence

$$t\mathcal{I} \simeq t\Delta^{op}$$

as the functors  $(\_)^\wedge$  and  $(\_)^\vee$  send cropped trees of positive degree to cropped trees of positive degree.

# an object of $t\mathcal{I}_+$








e.g. an object of  $t\mathcal{I}_+$  of degree 3



# conclusion

We have the following equivalences

$$\begin{array}{ccc} \mathit{Disk}^{\mathit{op}} & \xrightarrow{\sim} & \Theta \\ \downarrow \sim & & \downarrow \sim \\ \mathit{iI}^{\mathit{op}} & \xrightarrow{\sim} & \mathit{i}\Delta \\ \downarrow \sim & & \downarrow \sim \\ \mathit{tI}^{\mathit{op}} & \xrightarrow{\sim} & \mathit{t}\Delta \end{array}$$

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